

# **Evolution of genome assembly methods: solutions and delusions**

Anton Bankevich

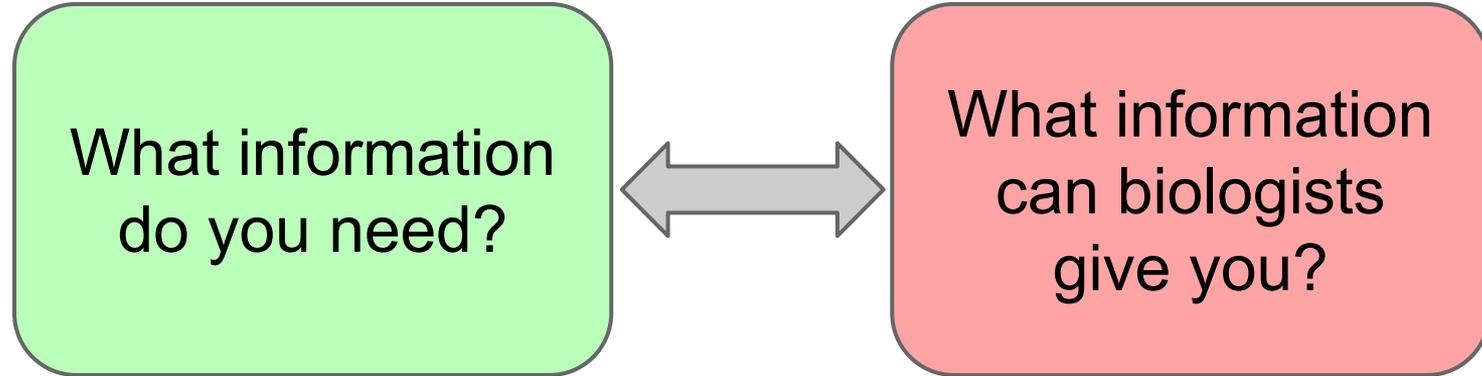
SPbAU 2014

# Assembly problem

- In 1944 the role of DNA as a primary information carrier was discovered.
- Before starting to understand information you should get it



# How can one reconstruct 3Gb long sequence?



Interaction between these two questions brings evolution into genome assembly problem.  
Along with one more small question.

# How can one reconstruct 3Gb long sequence?

## How much does it cost?

What information  
do you need?



What information  
can biologists  
give you?

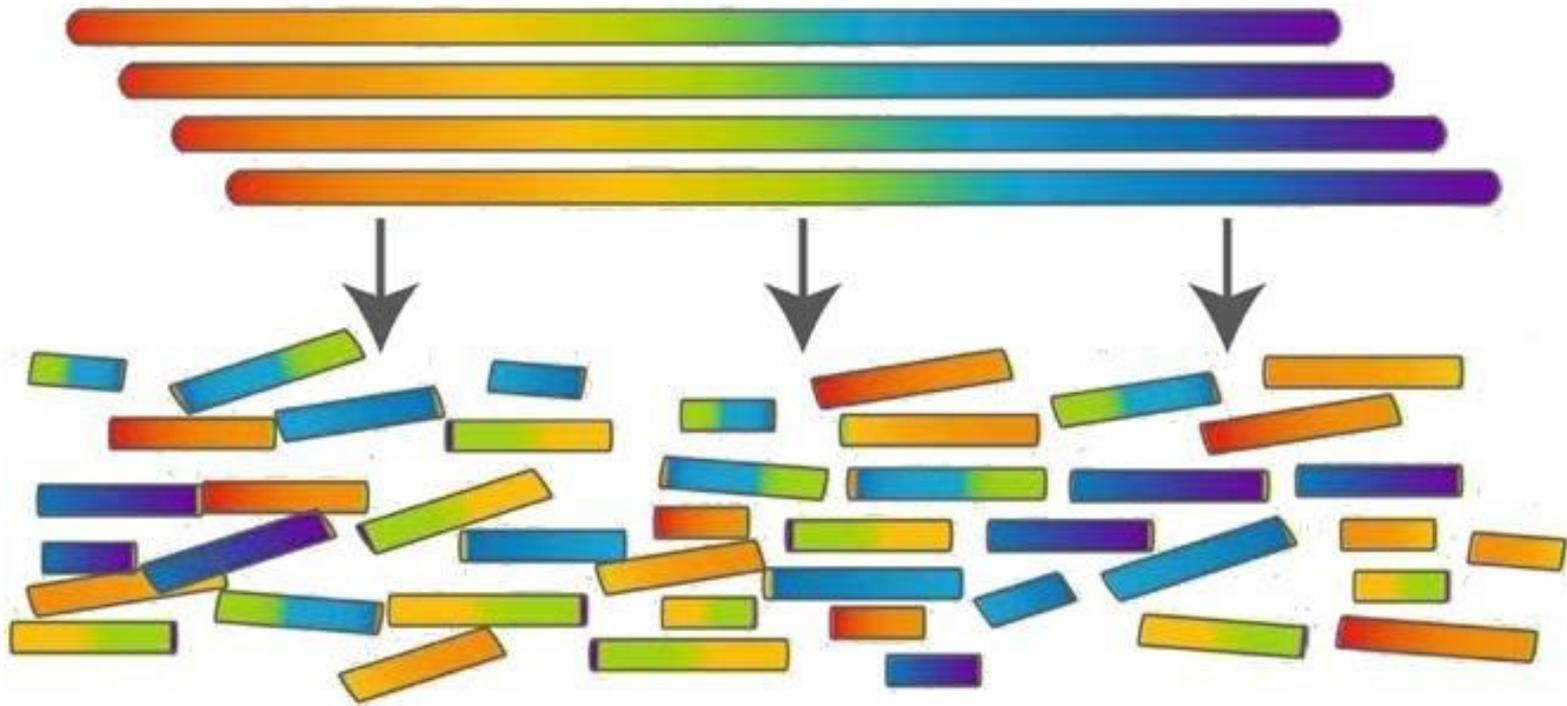
# Read concept

Sequences of small pieces DNA can be read.

Reconstruction of the whole genome turns into solution of a puzzle.



# Read concept



# Assembly problem

Given: a set of subsequences of genome sequence

Goal: reconstruct whole genome

Question: is it even possible?

# Shortest superstring delusion

The simplest  
solution is always  
right



Genome is a  
shortest  
superstring of all  
reads

# Shortest superstring delusion

~~The simplest solution is always right~~

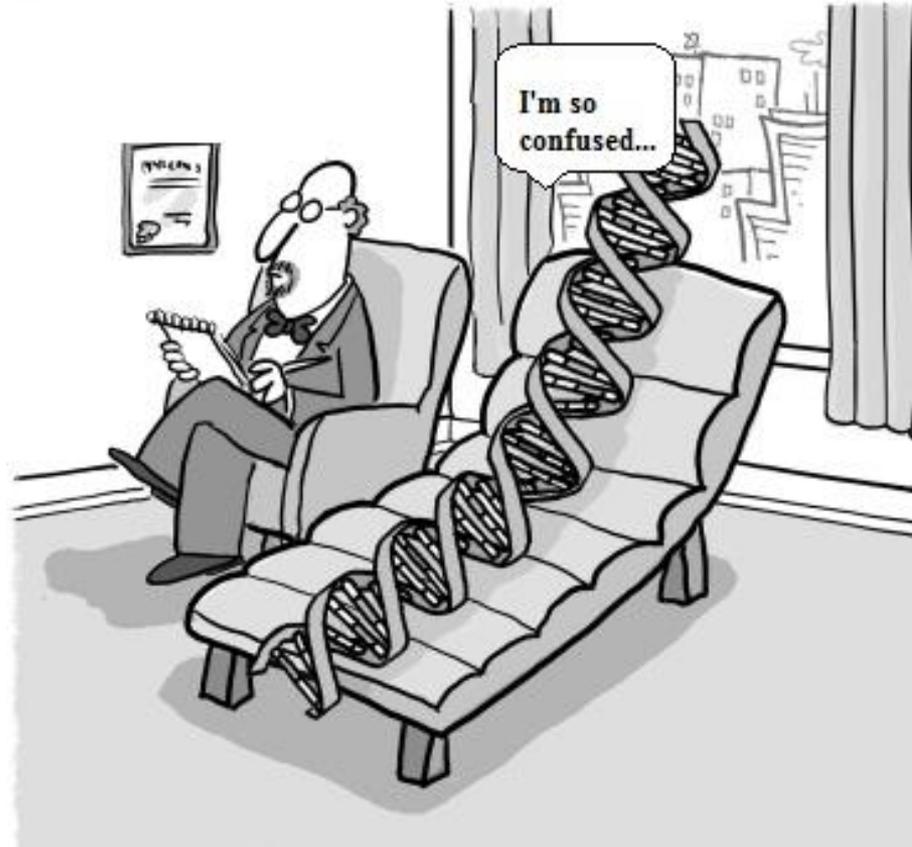


Genome is a shortest superstring of all reads

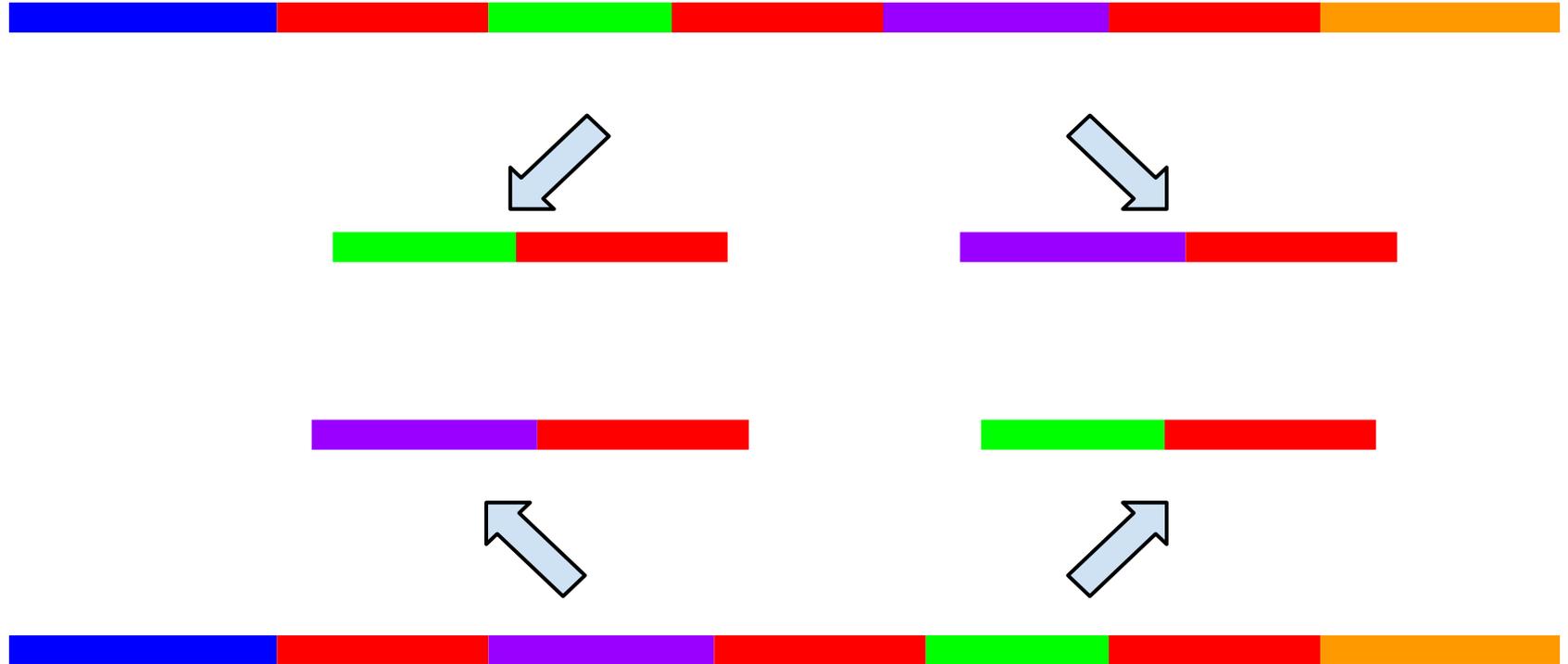
# Repeat structure

**Alu**: длина 300,  
кратность 1 000 000

**L1**: длина 6000,  
кратность 516 000



# Repeat structure



# Assembly problem

Formal statement:

# Assembly problem

Formal statement:

Has never been formulated...

# Assembly problem

## Wishes to assembly result:

- Answer is a set of sequences (contigs)
- We want longer contigs
- We want contigs to be subsequences of genome
- We do not want contigs to intersect

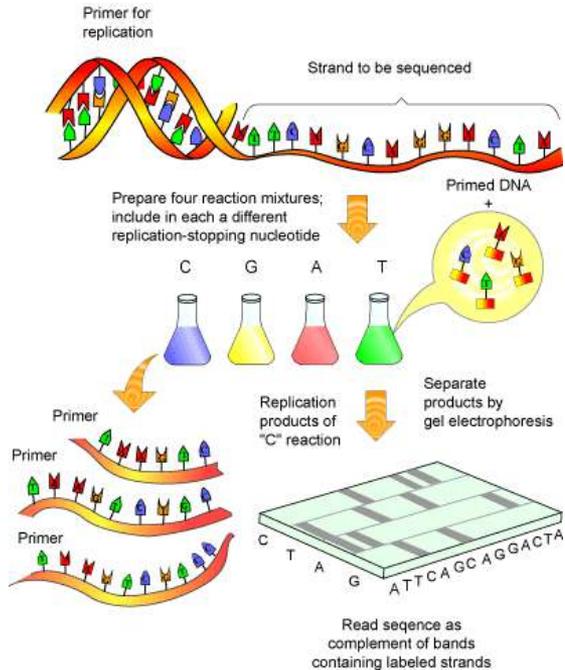
# Assembly problem

## Wishes to assembly result:

- Answer is a set of sequences (contigs)
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So what's the input?

# Sanger



Length  
Accuracy \*  
Cost per  $10^6$  bp

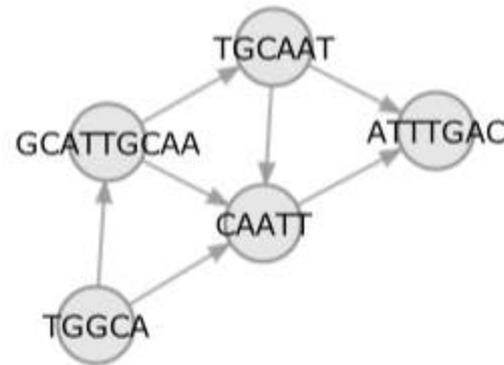
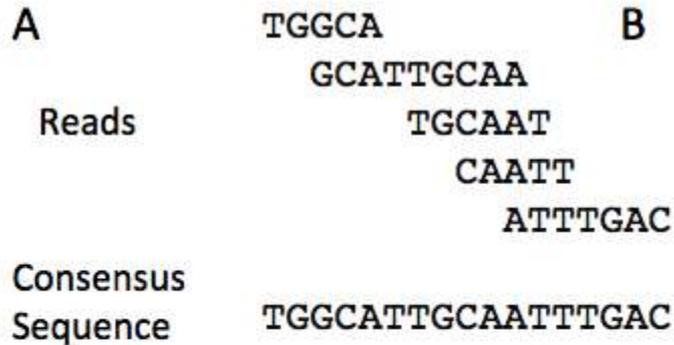
up to 900 bp  
99.9 %  
2400 \$

# Overlap graph: theory

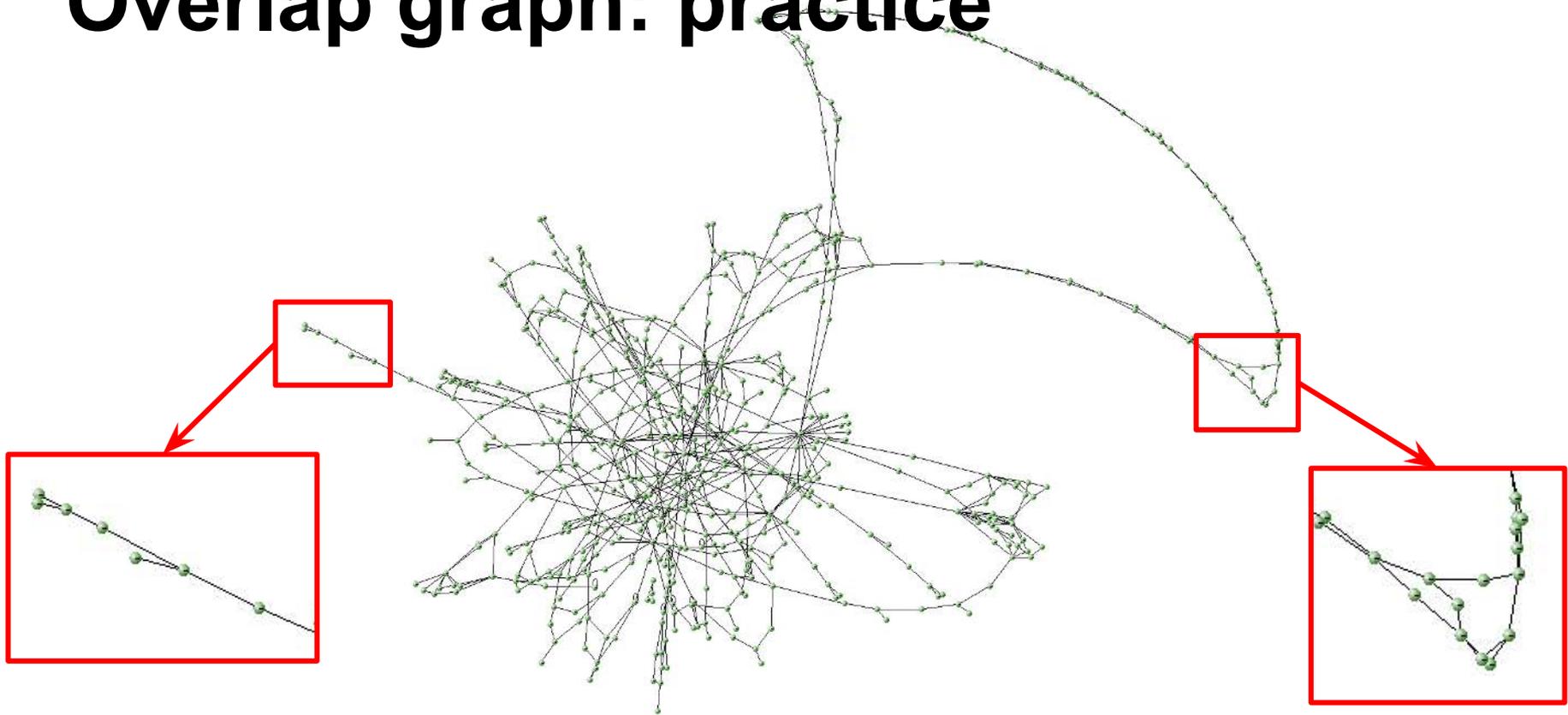
$R = \{r_1, \dots, r_n\}$  is a set of reads

$V = R$ ,  $e = (r_i, r_j)$  in  $E$ , if read  $r_i$  overlaps read  $r_j$

$w(e)$  - weight of edge  $e$  is a size of overlap



# Overlap graph: practice



overlap graph constructed from long reads for bacterial genome (output of Newbler)

# Delusions about overlap graph

- Finding Hamiltonian path of maximal weight **does not** give us superstring of  $R$  for genome with repeats
- Overlap graph constructed from reads contains erroneous edges due to sequencing errors
- $|V| = |R|$ ,  $\max |E| = R^2$

# Next Generation Sequencing

Starting from 2005 several new sequencing technologies hit the market: 454, Solexa (Illumina).

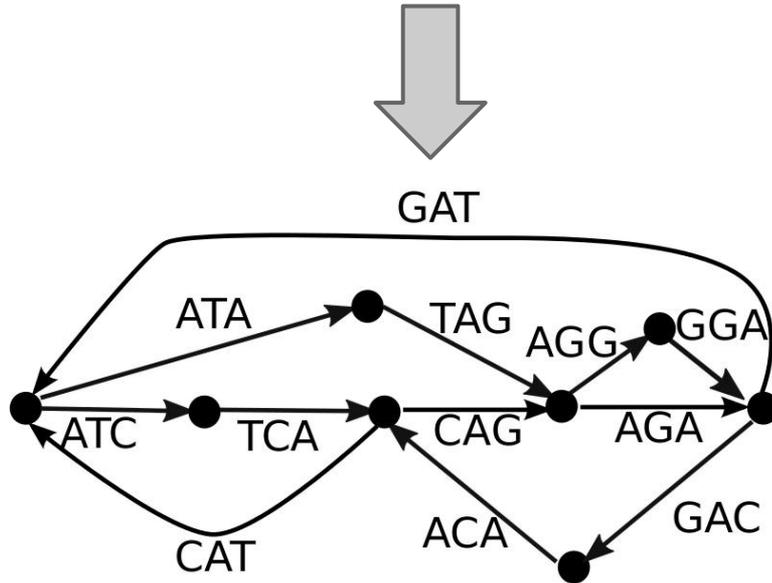
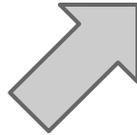
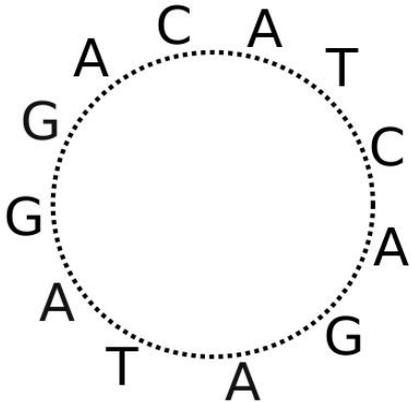
New technologies produced millions of short reads making overlap graph approach inapplicable.

# de Bruijn graph

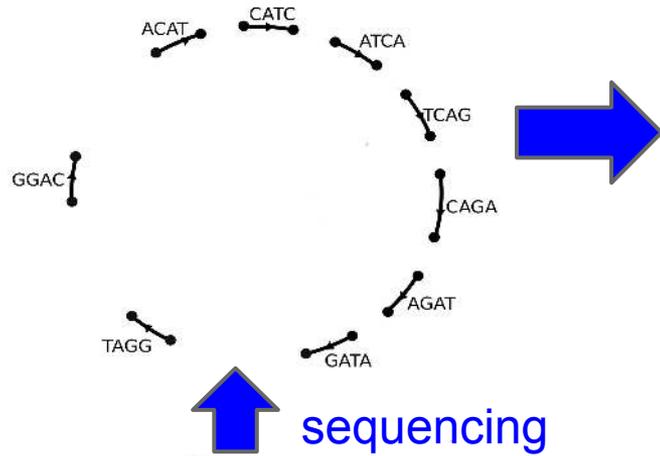
Vertices: k-mers from genome

Edges: (k+1)-mers from genome

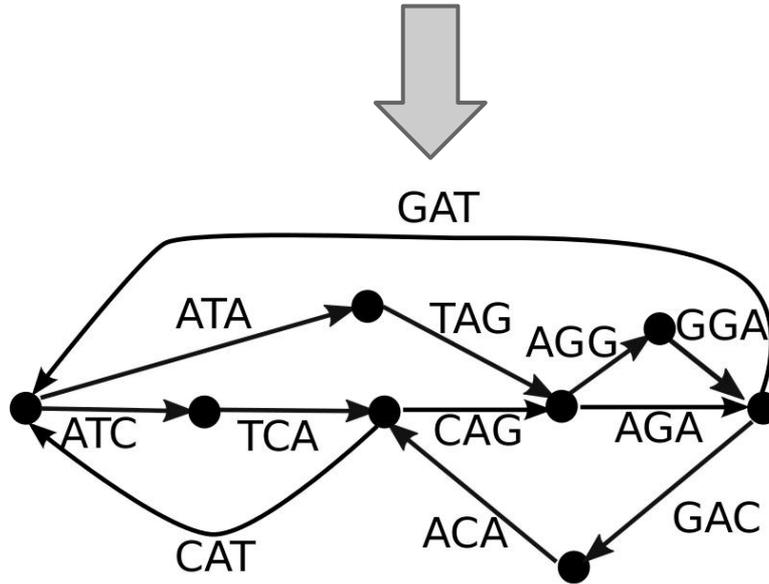
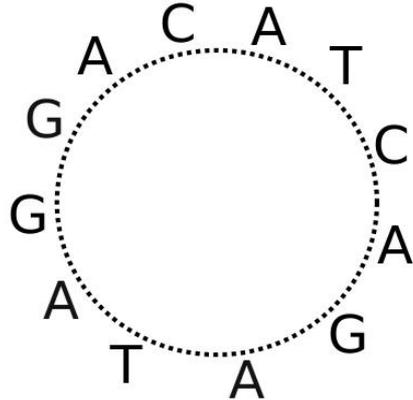
k=2: 3-mer ACG gives AC -> CG



# de Bruijn graph



Vertices: k-mers from reads  
Edges: (k+1)-mers from reads  
k=2: 3-mer ACG gives AC -> CG

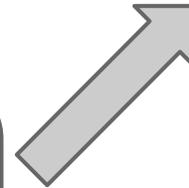
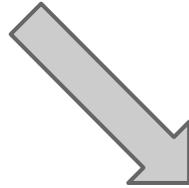


# Delusions about de Bruijn graph

Genome goes through all the edges of de Bruijn graph

Genome assembly is simple with de Bruijn graph since it is easy to find Eulerian cycle

Genome corresponds to Eulerian path in the graph

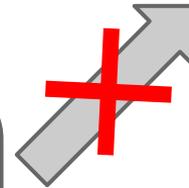
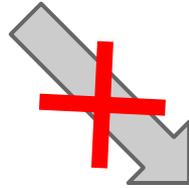


# Delusions about de Bruijn graph

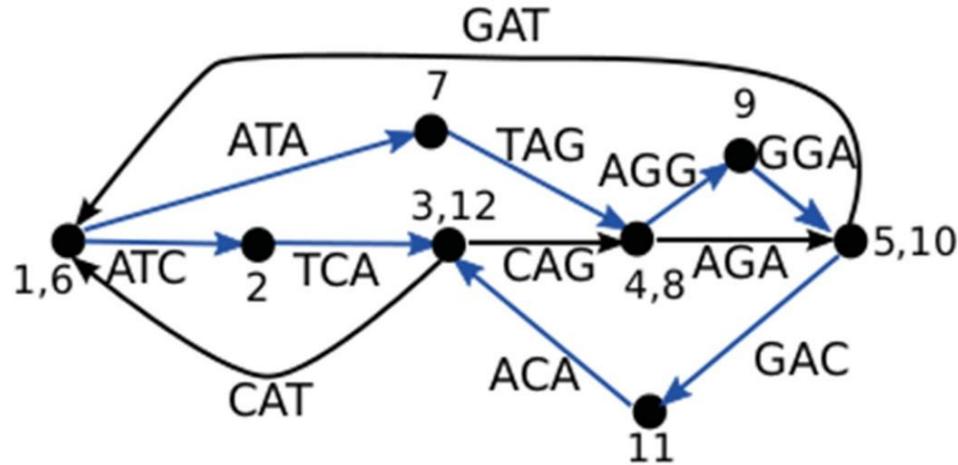
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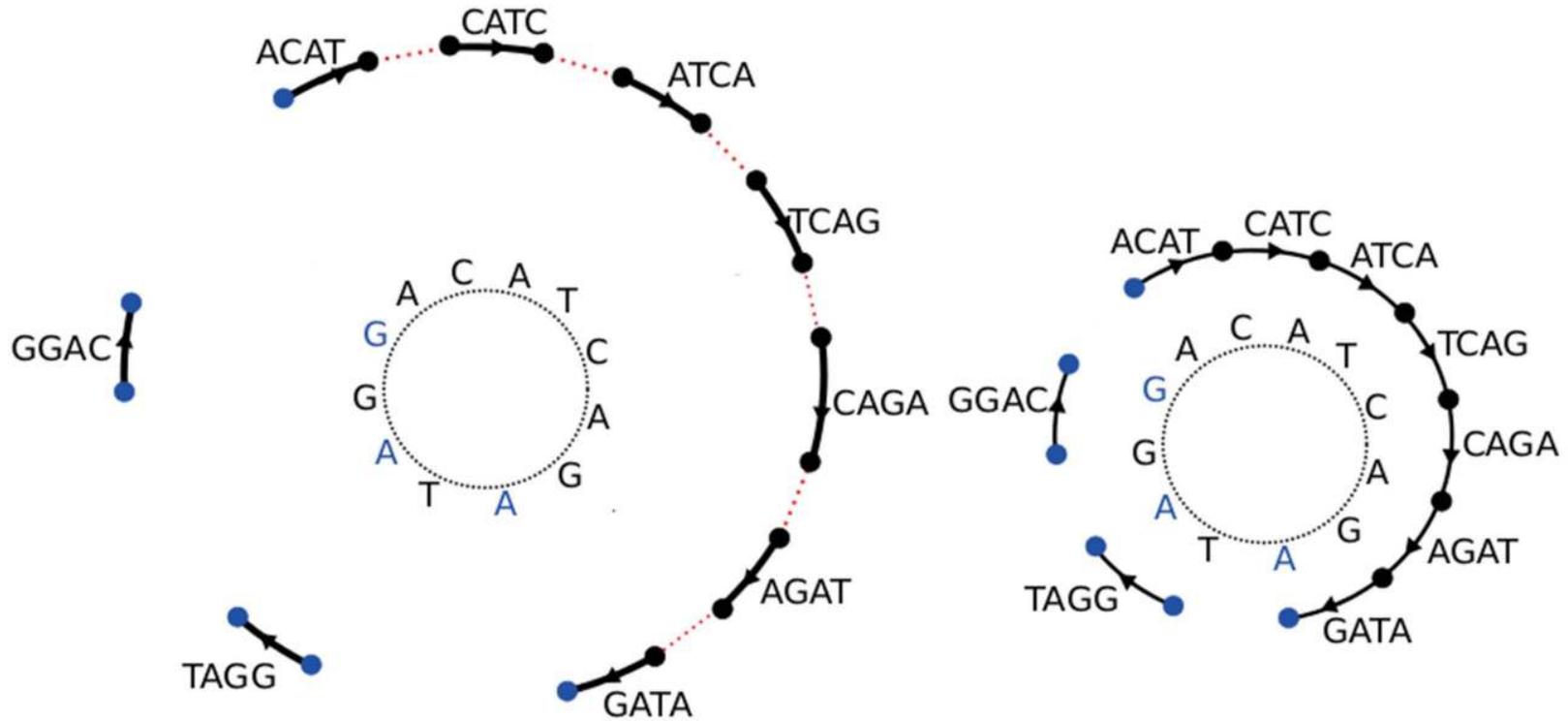
# Assembly using de Bruijn graph



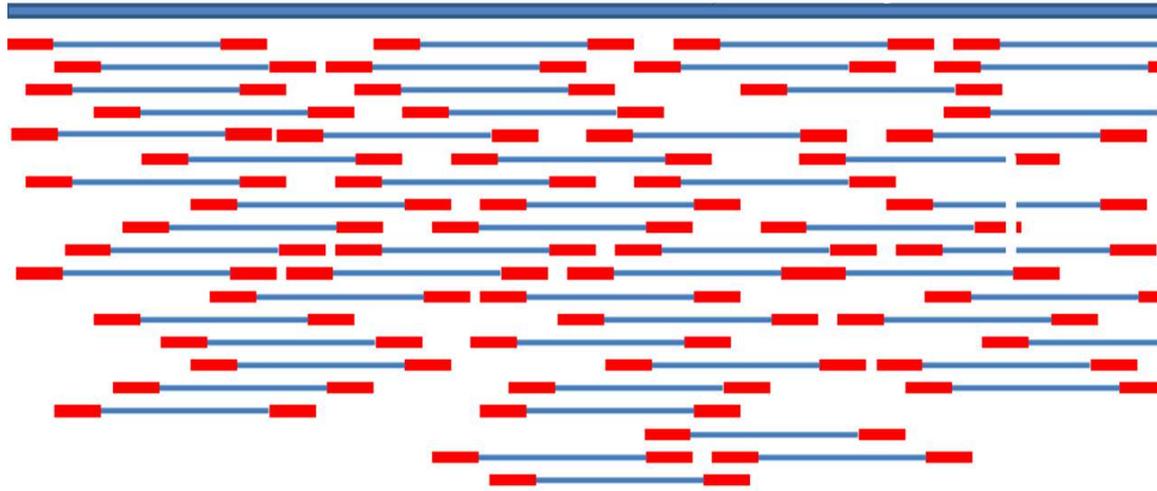
Output contigs are formed by non branching paths in de Bruijn graph (highlighted in blue)

Only repeats of length  $k$  can be resolved this way

# Choice of k



# Paired reads



Paired reads allow us to construct long genomic paths in de Bruijn graph instead of just using non branching paths.

# De Bruijn graph challenges

- De Bruijn graph of human genome contains 3 billion vertices.
- De Bruijn graph of reads from human genome contains at least 10 times more vertices.
- Even if we use 10 bytes/kmer we need 300 Gb memory
- If we consider all additional information and overheads we hit the roof.

# De Bruijn graph as a data structure

Operations of de Bruijn graph data structure

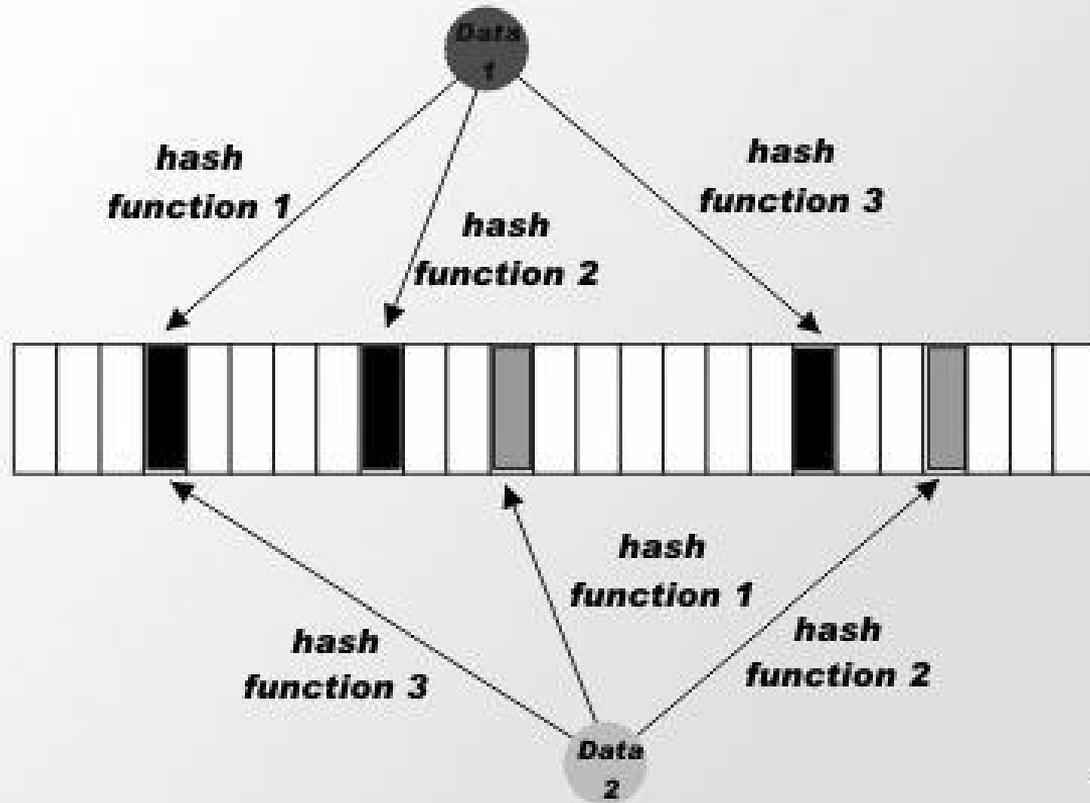
- Find all adjacent k-mers for a given k-mer
- Iterate over all k-mers

How to do it using as little memory as possible.

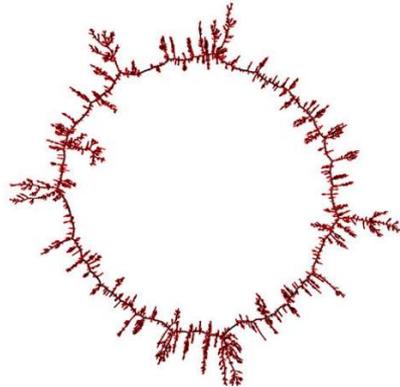
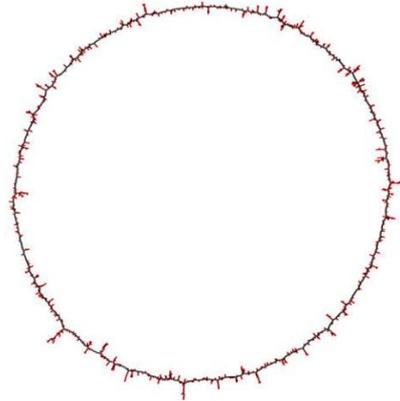
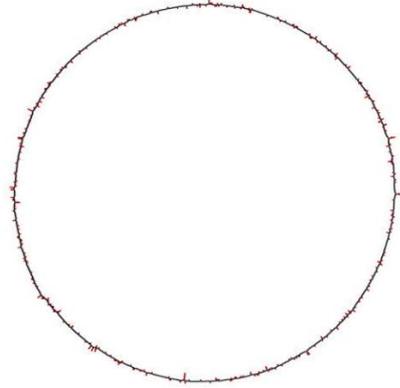
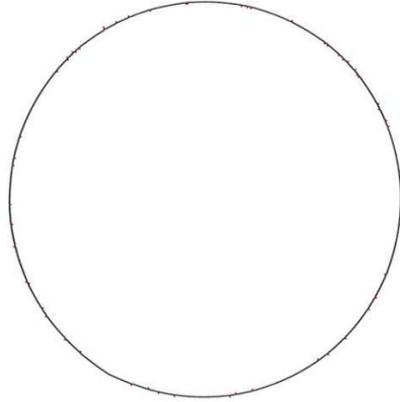
# Distributed storage (Abyss)

- Each k-mer is stored in one of the nodes based on some hash.
- Hash is chosen wisely in such a way that adjacent vertices have high probability to be stored on the same node

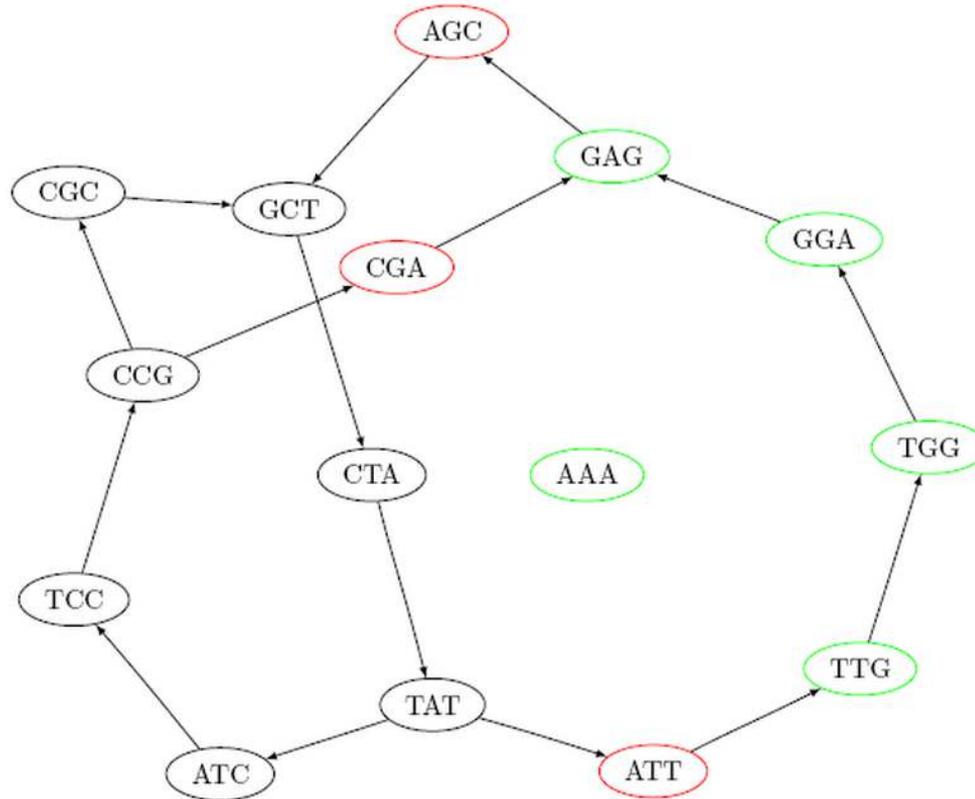
# Bloom filters (Minia)



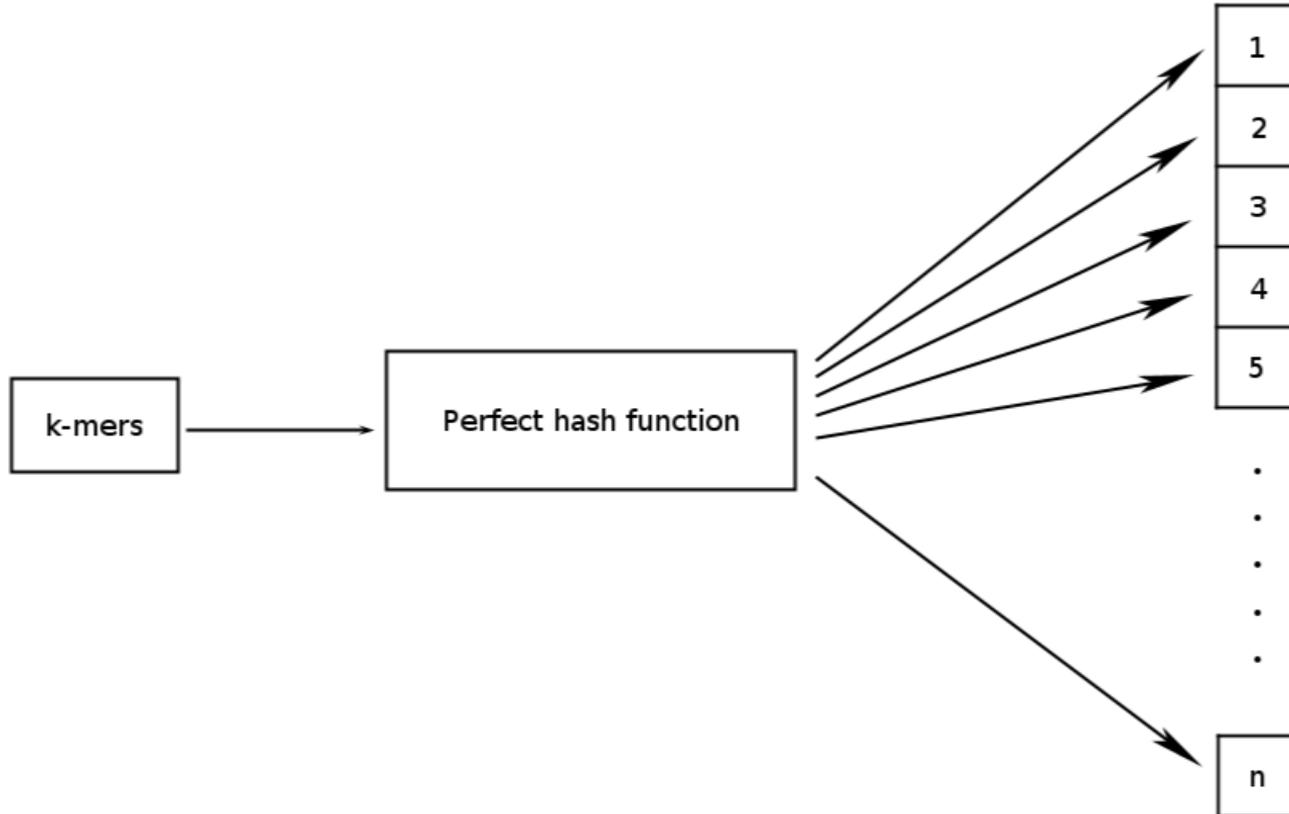
# Probabilistic de Bruijn graph



# Exact representation using bloom filters



# Perfect hashing (Miraculous, SPAdes)



# Perfect hashing

- Requires only 2.7 bits/kmer
- Any additional information can be stored without overhead
- Allows to modify de Bruijn graph
- Does not allow to check if given k-mer is in the graph

# **Sparse de Bruijn graphs (SparseAssembler)**

Idea: A lot of information in de Bruijn graph is not useful.

Solution: remove all highly covered and all low-covered k-mers/vertices and consider reads as edges.

Result: simultaneous graph cleaning and repeat resolution.

# Results of NGS era

- Assembly technology development result: Dozens of excellent assemblers appeared and used to assemble various species
- Sequencing technology development result: NGS technologies came close to the quality of Sanger. Today Illumina produces accurate and cheap reads with length up to 300bp.

# SGA

SGA assembler uses FM-index to construct overlap graph using very low amount of memory.

The problem of storing de Bruijn graph in FM-index remains open.

# Comparison of approaches

## Overlap graph

- Can use the whole read length for repeat resolution
- Is tolerant for moderate error rate in reads

## De Bruijn graph

- Can be used with very short reads (obsolete?)
- Can be used with read libraries of very high coverage
- Can be used with very high error rate

# Third generation and new challenges

NGS proved unable to assemble large and even complex small genomes efficiently.

Third generation is coming with more challenges:

- Pacbio technology provides very long but inaccurate reads
- Bionano technology creates maps of 100Kb long genome fragments
- TruSeq technology from Illumina provides 10Kb long, accurate but “virtual” reads.

Thank you

Questions?

# Reference assisted vs de novo

# Alignment problem

# MicroChips

# Genome map concept

# Technologies